

Classical Applications: A Summary

*WARNING: this is just a summary of the material covered in the full slide-deck **Classical Applications** that will orient you as per the topics covered there; you are required to learn the full version, not just this summary!*

- Optimal Product Mix and the Resource Allocation problem
- Generalized assignment problems: continuous case, machine scheduling
- Optimal portfolio: detecting arbitrage opportunities (optional)
- Make sure you study and solve all the examples in the full-slide deck!

Optimal Product Mix

- Linear programming permeates the entire fields of economic studies, management science, operations research, and logistics
- Perhaps the most prevalent example is the Product Mix Optimization (or Production Planning) problem
- A factory has finite resources available to manufacture goods/commodities
- Products can be sold to realize immediate profit or stocked in the hope for future profit
- Goal is to determine the optimal allocation of resources in order to maximize profit
- Assumptions:
 - products and market demands are independent
 - resources are divisible
 - prices and demands can be reliably predicted

Optimal Product Mix

- The quantity of goods produced in a period plus the stock must cover the demand in each period, subject to limits on resource availability
- Goal: maximize profit while minimizing storage costs
- Let the starting stock be $y_{i0} = 0$ for each i

$$\max \sum_{t=1}^T \sum_{i=1}^I c_{it} (y_{i,t-1} + x_{it} - y_{it}) - \sum_{t=1}^T \sum_{i=1}^I q_{it} y_{it}$$

$$\text{s.t. } y_{i,t-1} + x_{it} - y_{it} = d_{it} \quad i \in \{1, \dots, I\}, t \in \{1, \dots, T\}$$

$$\sum_{i=1}^I a_{ik} x_{it} \leq b_{kt} \quad k \in \{1, \dots, K\}, t \in \{1, \dots, T\}$$

$$x_{it}, y_{it} \geq 0 \quad i \in \{1, \dots, I\}, t \in \{1, \dots, T\}$$

Generalized Assignment Problem

- Another crucial application of linear programming in operations research
- Given m agents and n jobs that must be assigned to agents
- Each task can be performed by each agent, but with different effectiveness
 - agent i can do job j with w_{ij} units of effort, meanwhile we realize p_{ij} profits
 - agent i has w_i units of working capacity
 - b_j units of job j must be done
- Goal is to maximize profits
- Suppose that jobs are arbitrarily divisible
- Otherwise we get an integer linear program: NP-hard

Generalized Assignment Problem

- Let x_{ij} denote the quantity of job j performed by agent i [units]

$$\max \sum_{i=1}^m \sum_{j=1}^n p_{ij} x_{ij}$$

$$\sum_{j=1}^n w_{ij} x_{ij} \leq w_i \quad i \in \{1, \dots, m\}$$

$$\sum_{i=1}^m x_{ij} = b_j \quad j \in \{1, \dots, n\}$$

$$x_{ij} \geq 0 \quad i \in \{1, \dots, m\}, \\ j \in \{1, \dots, n\}$$